

Prelude + NASA Langley PVS Libraries

by

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There is one thing better than having a powerful set of tools and skills to carry out a mechanized mathematical proof.

**Discover that someone else has already done the proof
for you!**

Prelude

- The prelude is a special file containing theories with predefined types, functions, datatypes etc.
- It also contains a lot of lemmas we can use in proofs.
- It is organized in many theories.
- Inspect it by typing M-x vpf – “view prelude file”.
- The prelude is available **without** IMPORTING it.
- Some of it just documents things that are handled automatically by the system
- But, most of it is very useful.

Scope of the Prelude

```
39: equalities [T: TYPE]: THEORY
111: quantifier_props [t: TYPE]: THEORY
144: defined_types [t: TYPE]: THEORY
224: functions [D, R: TYPE]: THEORY
396: identity [T: TYPE]: THEORY
417: relations [T: TYPE]: THEORY
453: orders [T: TYPE]: THEORY
649: wf_induction [T: TYPE, <: (well_founded?[T])]: THEORY
666: measure_induction [T,M:TYPE, m:[T->M], <: (well_founded?)]
703: sets [T: TYPE]: THEORY
1199: function_inverse[D: NONEMPTY_TYPE, R: TYPE]: THEORY
1656: numbers: THEORY
1670: number_fields: THEORY
1730: reals: THEORY
1805: bounded_real_defs: THEORY
```

Scope of the Prelude (cont.)

```
1951: rationals: THEORY
2035: integers: THEORY
2978: exponentiation: THEORY
3242: divides : THEORY
3344: modulo_arithmetic : THEORY
3490: subrange inductions[i: int, j: upfrom(i)]: THEORY
3578: int_types[m: int] : THEORY
3588: nat_types[m: nat] : THEORY
3674: finite_sets[T: TYPE]: THEORY
3981: sequences[T: TYPE]: THEORY
4054: finite_sequences [T: TYPE]: THEORY
4217: list_props [T: TYPE]: THEORY
4698: bv[N: nat]: THEORY
4989: infinite_sets_def[T: TYPE]: THEORY
```

Functions

The theory functions provides the basic properties of functions

$f : D \rightarrow R$:

```
functions [D, R: TYPE]: THEORY
```

```
BEGIN
```

```
  f, g: VAR [D -> R]
```

```
  x, x1, x2: VAR D
```

```
  y: VAR R
```

```
extensionality_postulate: POSTULATE (FORALL (x:D): f(x) = g(x))  
                           IFF f = g
```

```
eta: LEMMA (LAMBDA (x: D): f(x)) = f
```

```
injective?(f): bool = (FORALL x1,x2: (f(x1)=f(x2) => (x1=x2)))
```

```
surjective?(f): bool = (FORALL y: (EXISTS x: f(x) = y))
```

```
bijection?(f): bool = injective?(f) & surjective?(f)
```

```
bij_is_inj: JUDGEMENT (bijection?) SUBTYPE_OF (injective?)
```

```
bij_is_surj: JUDGEMENT (bijection?) SUBTYPE_OF (surjective?)
```

Properties of numbers and operations

The theories (such as real_axiom, real_types, rationals and real_props) gives lemmas about the usual operations on numbers.

For example:

- Commutativity, associativity of operators (Decision procedures do this)
- Preservation of subtypes under certain operations (eg. $px * py$ is positive)
- Odd/even integers
- Cancellation laws to help prove non-linear equations and inequalities

Now that we have **manip** and **field** strategies you do not need to access these directly very often anymore.

orders

```
orders [T: TYPE]: THEORY
  x, y: VAR T
  <=, < : VAR pred[[T, T]]

  preorder?(<=): bool = reflexive?(<=) & transitive?(<=)
  partial_order?(<=): bool = preorder?(<=) & antisymmetric?(<=)
  strict_order?(<): bool = irreflexive?(<) & transitive?(<)
  total_order?(<=): bool = partial_order?(<=) & dichotomous?(<=)
  trichotomous?(<): bool = (FORALL x, y: x < y OR y < x OR x = y)

  upper_bound?(<)(x, pe): bool = FORALL (y: (pe)): y < x
  upper_bound?(<)(pe)(x): bool = upper_bound?(<)(x, pe)
  lower_bound?(<)(x, pe): bool = FORALL (y: (pe)): x < y
  lower_bound?(<)(pe)(x): bool = lower_bound?(<)(x, pe)

  least_upper_bound?(<)(x, pe): bool =
    upper_bound?(<)(x, pe) AND
    FORALL y: upper_bound?(<)(y, pe) IMPLIES (x < y OR x = y)

  greatest_lower_bound?(<)(x, pe): bool =
    lower_bound?(<)(x, pe) AND
    FORALL y: lower_bound?(<)(y, pe) IMPLIES (y < x OR x = y)
```

Standard Functions Provided by Prelude

abs(x)	– absolute value of <i>x</i>
even?(i), odd?(i)	– predicates over integers: even/odd
floor(x), ceiling(x)	– returns integer near <i>x</i> .
<i>x</i> ⁱ	– <i>x</i> to an integer power
rem(b)(x)	– remainder of <i>x/b</i>
exp2(n)	– 2^n
glb(S)	– greatest lower bound on set of reals
lub(S)	– least upper bound on set of reals
min(S)	– minimum of a set of natural numbers
sgn(m)	– sign of a real number
min(m,n)	– smallest of <i>m</i> and <i>n</i>
max(m,n)	– largest of <i>m</i> and <i>n</i>
ndiv(m,n)	– integer division <i>m/n</i> ¹
inverse(f)	– inverse of a function

set operators

sequence operators

bitvector operations

We can use libraries as bundles of theories covering a specific application area.

With the PVS distribution there are two libraries: bitvectors and finite_sets. Both contain strategies and can be found in your PVS/lib directory (same place as the prelude file) ².

There is also a collection of libraries available from NASA LaRC at
<http://shemesh.larc.nasa.gov/fm/ftp/larc/PVS-library/pvslib.html>

²subdirectory of your PVS installation directory (i.e. PVSPATH)

Installation of NASA libraries

Download the NASA library gzipped tar file to the main directory (PVSPATH). Then execute

```
tar xvfz p50_pvslib-full.tgz
```

This will create a subdirectory nasalib. You must set the environment variable PVS_LIBRARY_PATH to the nasalib directory. In C shell (csh, tcsh, etc):

```
setenv PVS_LIBRARY_PATH "<pvs-dir>/nasalib"
```

In Bourne shell (sh, bash, etc):

```
export PVS_LIBRARY_PATH=<pvs-dir>/nasalib"
```

Then, add the following line to the file ~/.pvs.lisp (create it if it doesn't exist):

```
(load "<pvsdir>/nasalib/pvs-patches.lisp")
```

We strongly recommend recreating the binary files if you have installed the no-binary distribution. This can be accomplished by issuing the following command from the directory <pvsdir>/nasalib:

```
../provethem nasalib.all
```

IMPORTING Library Theories

If you have set the PVS_LIBRARY_PATH environment variable correctly, then the following statements in your specifications

```
IMPORTING reals@sigma, vectors@closest_approach_2D
```

will direct the PVS system will look for subdirectories `reals` and `vectors` to find the `sigma` and `closest_approach_2D` theories.

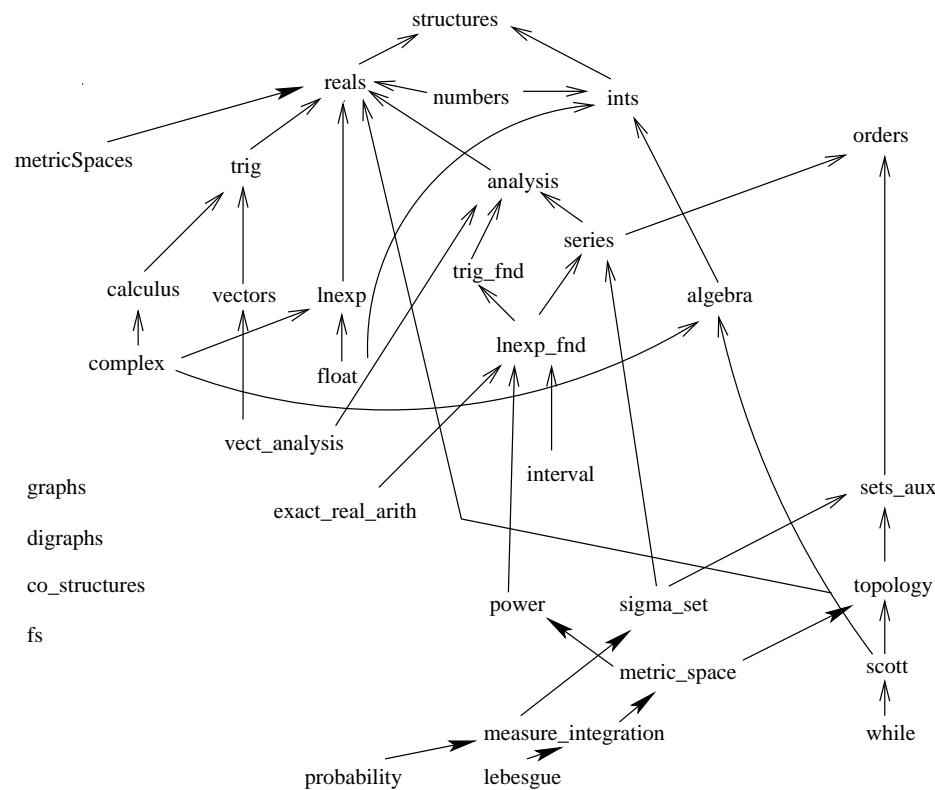
The LIBRARY statement is still available though its use is **deprecated**.

```
realslib: LIBRARY = "/home/rwb/lib/reals "
```

```
IMPORTING realslib@quadratic
```

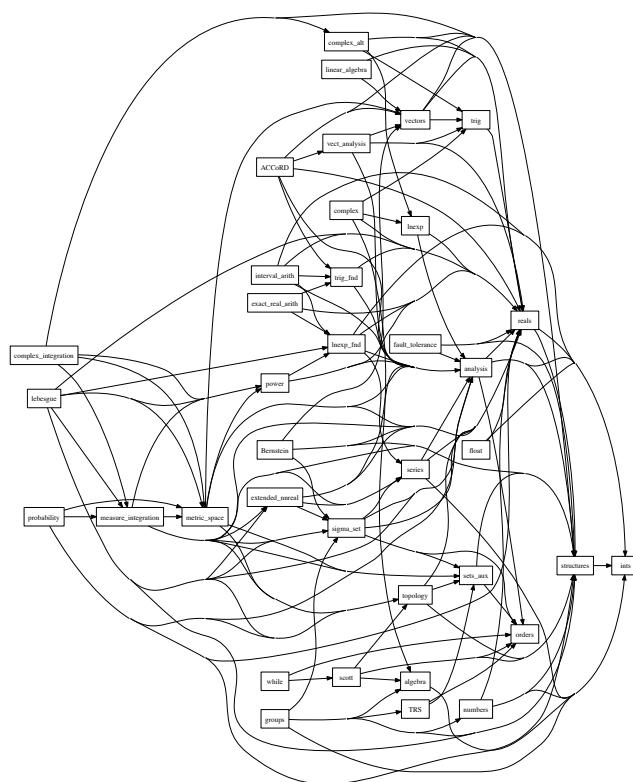
This path variable PVS_LIBRARY_PATH which tells PVS where to look is a **list**, so libraries can be stored in multiple places, however, we favor placing them all in one directory.

Structure of NASA libraries



Structure of NASA libraries (cont.)

- The picture above does not include the following newer libraries: fault_tolerance, TRS, groups, interval_arith, ACCoRD, algexp, Bernstein, complex_integration,
- There are over 12,000 theorems available.
- You can use Hypatheon to help you find what you need.



The NASA reals Library

```
top_sigma      -- provides a family of summation functions
product_real   -- defines product over functions [int -> real]
bounded_reals  -- defines sup, inf, max, min
real_sets       -- properties of sup, inf, max, min
real_fun_ops,   -- adding, subtracting, etc on functions
real_fun_preds -- increasing?, decreasing?, etc on functions
real_fun_props -- properties about defs in real_fun_preds
abs_lems        -- additional properties of abs
exponent_props -- additional properties of expt
sqrt            -- properties of square root
sqrt_approx     -- definitions of sqrt_newton and sqrt_bisect
                  (newton +bisection methods for computing sq. root)
sq               -- square function and properties
sq_rew           -- installs useful AUTO_REWRITE+s
sign             -- properties of sign function
quadratic       -- solution of quadratic equations
quadratic_minmax -- Minimum and Maximum of quadratic equations
quadratic_ax    -- solution of quadratic equations (Obsolete)
binomial,        -- Binomial coefficient
polynomials      -- Polynomials
```

The sqrt Theory in the reals Library

```
sqrt(nnx): {nnz | nnz*nnz = nnx}
sq(a): nonneg_real = a*a

sqrt_def      : LEMMA sqrt(nnx) * sqrt(nnx) = nnx
sqrt_square   : LEMMA sqrt(nnx * nnx) = nnx
sqrt_sq       : LEMMA x >= 0 IMPLIES sqrt(sq(x)) = x
sqrt_sq_neg   : LEMMA x < 0 IMPLIES sqrt(sq(x)) = -x
sqrt_sq_abs   : LEMMA sqrt(sq(x)) = abs(x)
sqrt_times    : LEMMA sqrt(nny * nnz) = sqrt(nny) * sqrt(nnz)
sqrt_div      : LEMMA nnz /= 0 IMPLIES
                  sqrt(nny / nnz) = sqrt(nny) / sqrt(nnz)
```

% ----- Inequalities -----

```
sqrt_lt       : LEMMA sqrt(nny) < sqrt(nnz) IFF nny < nnz
sqrt_le       : LEMMA sqrt(nny) <= sqrt(nnz) IFF nny <= nnz
sqrt_gt       : LEMMA sqrt(nny) > sqrt(nnz) IFF nny > nnz
sqrt_ge       : LEMMA sqrt(nny) >= sqrt(nnz) IFF nny >= nnz
sqrt_eq       : LEMMA sqrt(nny) = sqrt(nnz) IFF nny = nnz
sqrt_less     : LEMMA nnx > 1 IMPLIES sqrt(nnx) < nnx
sqrt_more     : LEMMA nnx > 0 AND nnx < 1 IMPLIES sqrt(nnx) > nnx
```

Sigma Theories in reals Library

The sigma theory introduces and defines properties of the sigma function that sum an arbitrary function $F: [T \rightarrow \text{real}]$ over a range from low to high

$$\text{sigma}(\text{low}, \text{high}, F) = \sum_{j=\text{low}}^{\text{high}} F(j)$$

The most general “sigma” theory is sigma:

```
sigma[T: TYPE FROM int]: THEORY
  low,high: VAR T
  F: VAR function[T -> real]

  sigma(low, high, F): RECURSIVE real
    = IF low > high THEN 0
      ELSIF high = low THEN F(low)
      ELSE F(high) + sigma(low, high-1, F)
      ENDIF
  MEASURE (LAMBDA low, high, F: abs(high-low))
```

which says that the index type is a subtype of the integers³ and the function F is real-valued.

³T can be any subtype of the integers that is connected.

But remember that

```
F: [nat -> real]
F: [int -> real]
F: [posnat -> real]
F: [below[N] -> real]
F: [upto[N] -> real]
```

all have different domains, so there are special theories for each

```
sigma_nat,          % sigma over functions [nat --> real]
sigma_posnat,       % sigma over functions [posnat --> real]
sigma_int,          % sigma over functions [int --> real]
sigma_upsto,         % sigma over functions [upto[N] --> real]
sigma_below,         % sigma over functions [below[N] --> real]
sigma_below_sub,    % equality of sigmas with different domains
sigma,              % generic theory
```

Sigma cont.

NOTE: Summations over functions with an arbitrary number of arguments is accommodated by the following technique:

```
G(x,y,z: real,n: nat): real  
IMPORTING sigma[nat]      % or sigma_nat  
sum = sigma(low,high, (LAMBDA (n:nat): G(x,y,z,n))
```

Integers (ints) Library

```
div                  -- Ada Reference Manual division theory  
div_alt              -- Ada div defined recursively  
rem                  -- Ada Reference Manual rem theory  
mod_div_lems         -- relates mod to rem
```

This div truncates toward zero on a negative argument.

```
mod                  -- mod theory  
mod_lems             -- Additional lemmas about mod
```

Other theories:

```
max_upto,            -- max of a set of upto  
max_below,           -- max of a set of below  
div_nat,              -- integer division over nats  
mod_nat,              -- mod over nats  
abstract_min,         -- defines min over type T satisfying P  
abstract_max          -- defines max over type T satisfying P
```

See number_theory library for a div that truncates away from zero on a negative argument:

```
div_nt                -- Number theory division theory  
div_nt_alt            -- Number theory div defined recursively
```

The Trigonometry (trig) Library

There are actually two libraries:

- the **foundational** version: IMPORTING trig_fnd@trig_basic
- the **axiomatic** version: IMPORTING trig@trig_basic
 - typechecks significantly faster
 - derived from the foundational version to minimize the possibility that we created any “false” axioms.

Much of the foundational development was based on

Abramovitz and Stegun, "Handbook of Mathematical Functions"
Dover. 9th Edition (1972).

The foundational approach used was to define atan as an integral:

```
atan(x) = Integral(0, x, (LAMBDA (x:real):1/(1+x*x))) % 4.4.3
```

Then

```
asin(x:real_abs_le1): real_abs_le_pi2 =
  IF x = -1 THEN -pi/2 ELSIF x < 1
  THEN atan(x/sqrt(1-x*x)) ELSE pi/2 ENDIF

acos(x:real_abs_le1):nnreal_le_pi = pi/2 - asin(x)
```

Trig Library (Axiomatic)

```
trig_range : TYPE = {a | -1 <= a AND a <= 1}

sin(x:real) : real % = sin_phase(x-2*pi*floor(x/(2*pi)))
cos(x:real) : real % = cos_phase(x-2*pi*floor(x/(2*pi)))

sin_range_ax: AXIOM -1 <= sin(a) AND sin(a) <= 1
cos_range_ax: AXIOM -1 <= cos(a) AND cos(a) <= 1

sin_range : JUDGEMENT sin(a) HAS_TYPE trig_range
cos_range : JUDGEMENT cos(a) HAS_TYPE trig_range

Tan?(a) : bool = cos(a) /= 0

tan(a:(Tan?)) : real = sin(a)/cos(a)
```

Trig Library Continued

% ----- Pythagorean Property -----

```
sin2_cos2      : AXIOM sq(sin(a)) + sq(cos(a)) = 1
cos2          : LEMMA sq(cos(a)) = 1 - sq(sin(a))
sin2          : LEMMA sq(sin(a)) = 1 - sq(cos(a))
```

% ----- Basic Values -----

```
cos_0          : AXIOM cos(0) = 1
sin_0          : LEMMA sin(0) = 0
sin_pi2        : AXIOM sin(pi/2) = 1
cos_pi2        : LEMMA cos(pi/2) = 0
tan_0          : LEMMA tan(0) = 0
```

% ----- Negative Properties -----

```
sin_neg : AXIOM sin(-a) = -sin(a)
cos_neg : AXIOM cos(-a) = cos(a)
tan_neg : LEMMA FORALL(a:(Tan?)): tan(-a) = -tan(a)
```

Trig Library Continued

% ----- Co-Function Identities -----

```
cos_sin       : AXIOM cos(a) = sin(a+pi/2)
sin_cos       : LEMMA sin(a) = -cos(a+pi/2)
sin_shift     : LEMMA sin(pi/2 - a) = cos(a)
cos_shift     : LEMMA cos(pi/2 - a) = sin(a)
```

% ----- Arguments involving pi -----

```
neg_cos       : LEMMA -cos(a) = cos(a+pi)
neg_sin       : LEMMA -sin(a) = sin(a+pi)
```

```
sin_pi        : AXIOM sin(pi) = 0
cos_pi        : AXIOM cos(pi) = -1
tan_pi        : LEMMA tan(pi) = 0
sin_3pi2      : LEMMA sin(3*pi/2) = -1
cos_3pi2      : LEMMA cos(3*pi/2) = 0
sin_2pi        : AXIOM sin(2*pi) = 0
cos_2pi        : AXIOM cos(2*pi) = 1
tan_2pi        : LEMMA tan(2*pi) = 0
```

Trig Library Continued

```
% ----- Sum and Difference of Two Angles -----  
  
sin_plus : AXIOM sin(a + b) = sin(a)*cos(b) + cos(a)*sin(b)  
sin_minus : LEMMA sin(a - b) = sin(a)*cos(b) - sin(b)*cos(a)  
  
cos_plus : LEMMA cos(a + b) = cos(a)*cos(b) - sin(a)*sin(b)  
cos_minus : LEMMA cos(a - b) = cos(a)*cos(b) + sin(a)*sin(b)  
  
tan_plus : LEMMA Tan?(a) AND Tan?(b) AND Tan?(a+b) AND  
           tan(a) * tan(b) /= 1 IMPLIES  
           tan(a + b) = (tan(a)+tan(b))/(1-tan(a)*tan(b))  
  
tan_minus : LEMMA Tan?(a) AND Tan?(b) AND Tan?(a-b) AND  
           tan(a) * tan(b) /= -1 IMPLIES  
           tan(a - b) = (tan(a)-tan(b))/(1+tan(a)*tan(b))  
  
arc_sin_cos : AXIOM sq(a)+sq(b)=sq(c) IMPLIES  
              EXISTS d: a = c*cos(d) AND b = c*sin(d)  
  
pythagorean : LEMMA sq(a)+sq(b)=sq(nnc) IMPLIES  
               EXISTS(al:real): nnc=a*cos(al)+b*sin(al)
```

```
% ----- Double Angle Formulas -----
```

```
sin_2a      : LEMMA sin(2*a) = 2 * sin(a) * cos(a)  
cos_2a      : LEMMA cos(2*a) = cos(a)*cos(a) - sin(a)*sin(a)  
cos_2a_cos : LEMMA cos(2*a) = 2 * cos(a)*cos(a) - 1  
cos_2a_sin : LEMMA cos(2*a) = 1 - 2 * sin(a)*sin(a)  
tan_2a      : LEMMA Tan?(a) AND Tan?(2*a) AND  
              tan(a) * tan(a) /= 1 IMPLIES  
              tan(2*a) = 2 * tan(a)/(1 - tan(a)*tan(a))
```

```
% ----- Characterization of zeroes -----
```

```
sin_eq_0     : AXIOM sin(a) = 0 IFF EXISTS i: a = i * pi  
cos_eq_0     : LEMMA cos(a) = 0 IFF EXISTS i: a = i * pi + pi/2  
sin_eq_0_2pi : COROLLARY 0 <= a AND a <= 2*pi IMPLIES  
              (sin(a)=0 IFF a=0 OR a=pi OR a=2*pi)  
cos_eq_0_2pi : COROLLARY 0 <= a AND a <= 2*pi IMPLIES  
              (cos(a)=0 IFF a=pi/2 OR a=3*pi/2)
```

Still More Trig Library Theories

```
trig_basic      -- basic properties
trig_values     -- values of functions for special arguments
trig_ineq       -- trig inequalities
trig_extra      -- sum and product half-angle reductions and zeros
trig_approx     -- taylor series approximations to trig functions:
tan_approx      -- approximations for tangent
law_cosines     -- law of cosines
trig_degree     -- conversions to degrees
trig_inverses   -- inverse functions
trig_rew        -- auto-rewrites
asin            -- asin properties
acos            -- acos properties
atan            -- atan properties
atan2           -- two-argument arc tangent
atan2_props    -- additional properties of atan2
```

Example Using Trig Library

```
lib_ex: THEORY
BEGIN

IMPORTING trig@trig_basic

v,p: VAR posreal
x: VAR nnreal

syst(v,p)(x): real = sin(pi/2 - x)* sin(2*x)/(2*sin(x))

sp3: LEMMA syst(v,p)(x) = 1 - sq(sin(x))

sp4: LEMMA syst(v,p) = (LAMBDA x: 1 - sq(sin(x)))

END lib_ex
```

M-x pr <sp3>

Example Using Trig Library

sp3 :

```
|-----  
{1} FORALL (p, v: posreal, x: nnreal): syst(v, p)(x)  
= 1 - sq(sin(x))
```

Rule? (**skosimp***)

Repeatedly Skolemizing and flattening,
this simplifies to:

sp3 :

```
|-----  
{1} syst(v!1, p!1)(x!1) = 1 - sq(sin(x!1))
```

Rule? (**expand "syst"**)

Expanding the definition of syst,
this simplifies to:

sp3 :

```
|-----  
{1} sin(pi / 2 - x!1) * sin(2 * x!1) / (2 * sin(x!1))  
= 1 - sq(sin(x!1))
```

Example Using Trig Library (cont.)

Rule? (**lemma "sin_shift"**)

Applying sin_shift

this simplifies to:

sp3 :

```
{-1} FORALL (a: real): sin(pi / 2 - a) = cos(a)  
|-----  
[1] sin(pi / 2 - x!1) * sin(2 * x!1) / (2 * sin(x!1))  
= 1 - sq(sin(x!1))
```

Rule? (**inst?**)

Found substitution: a: real gets x!1,

Using template: sin(pi / 2 - a)

Instantiating quantified variables, this simplifies to:

sp3 :

```
{-1} sin(pi / 2 - x!1) = cos(x!1)  
|-----  
[1] sin(pi / 2 - x!1) * sin(2 * x!1) / (2 * sin(x!1)) = 1 - sq(sin
```

Example Using Trig Library (cont.)

sp3 :

```
{-1} sin(pi / 2 - x!1) = cos(x!1)
|-----
[1] sin(pi / 2 - x!1) * sin(2 * x!1) / (2 * sin(x!1)) = 1 - sq(sin
```

Rule? (replace -1)

Replacing using formula -1,

this simplifies to:

sp3 :

```
[-1] sin(pi / 2 - x!1) = cos(x!1)
|-----
{1} cos(x!1) * sin(2 * x!1) / (2 * sin(x!1)) = 1 - sq(sin(x!1))
```

Rule? (hide -1)

Hiding formulas: -1,

this simplifies to:

sp3 :

```
|-----
[1] cos(x!1) * sin(2 * x!1) / (2 * sin(x!1)) = 1 - sq(sin(x!1))
```

Example Using Trig Library (cont.)

```
|-----
[1] cos(x!1) * sin(2 * x!1) / (2 * sin(x!1)) = 1 - sq(sin(x!1))
```

- Let's use sin double angle formula:

sin_2a : LEMMA $\sin(2a) = 2 * \sin(a) * \cos(a)$

- Will use it as a rewrite rule

Rule? (rewrite "sin_2a")

Found matching substitution:

a: real gets x!1,

Rewriting using sin_2a, matching in *,
this simplifies to:

sp3 :

```
|-----
{1} cos(x!1) * (2 * (cos(x!1) * sin(x!1))) / (2 * sin(x!1)) =
    1 - sq(sin(x!1))
```

Example Using Trig Library (cont.)

sp3 :

$$\{1\} \frac{\cos(x!1) * (2 * (\cos(x!1) * \sin(x!1))) / (2 * \sin(x!1))}{1 - \text{sq}(\sin(x!1))}$$

Rule? (field 1)

Simplifying formula 1 with FIELD, this simplifies to:

sp3 :

$$\{1\} \frac{(\cos(x!1) * \cos(x!1))}{1 - \text{sq}(\sin(x!1))}$$

Example Using Trig Library (cont.)

sp3 :

$$\{1\} \frac{(\cos(x!1) * \cos(x!1))}{1 - \text{sq}(\sin(x!1))}$$

Rule? (lemma "sin2_cos2")

Applying sin2_cos2

this simplifies to:

$$\begin{aligned} \{-1\} \quad & \text{FORALL } (a: \text{real}): \text{sq}(\sin(a)) + \text{sq}(\cos(a)) = 1 \\ [1] \quad & (\cos(x!1) * \cos(x!1)) = 1 - \text{sq}(\sin(x!1)) \end{aligned}$$

Rule? (inst?)

Found substitution: a: real gets x!1, Using template: sq(sin(a))

Instantiating quantified variables, this simplifies to:

$$\begin{aligned} \{-1\} \quad & \text{sq}(\sin(x!1)) + \text{sq}(\cos(x!1)) = 1 \\ [1] \quad & (\cos(x!1) * \cos(x!1)) = 1 - \text{sq}(\sin(x!1)) \end{aligned}$$

Example Using Trig Library (cont.)

```
{-1} sq(sin(x!1)) + sq(cos(x!1)) = 1
|-----
[1] (cos(x!1) * cos(x!1)) = 1 - sq(sin(x!1))
```

Rule? (assert)

Simplifying, rewriting, and recording with decision procedures, this

```
{-1} sq(cos(x!1)) + sq(sin(x!1)) = 1
|-----
[1] (cos(x!1) * cos(x!1)) = 1 - sq(sin(x!1))
```

- No change

- Why?

Example Using Trig Library (cont.)

sp3 :

```
{-1} sq(cos(x!1)) + sq(sin(x!1)) = 1
|-----
[1] cos(x!1) * cos(x!1) = 1 - sq(sin(x!1))
```

Rule? (rewrite "sq_rew")

Found matching substitution:a: real gets cos(x!1),
Rewriting using sq_rew, matching in *, this simplifies to:
sp3 :

```
[1] sq(cos(x!1)) + sq(sin(x!1)) = 1
|-----
{1} sq(cos(x!1)) = 1 - sq(sin(x!1))
```

Rule? (assert)

Simplifying, rewriting, and recording with decision procedures,
Q.E.D.

Example Using Trig Library (cont.)

```
lib_ex: THEORY
BEGIN

  IMPORTING trig@trig_basic

  v,p: VAR posreal
  x: VAR nnreal

  syst(v,p)(x): real = sin(pi/2 - x)* sin(2*x)/(2*sin(x))

  sp3: LEMMA syst(v,p)(x) = 1 - sq(sin(x))

  sp4: LEMMA syst(v,p) = (LAMBDA x: 1 - sq(sin(x)))

END lib_ex
```

- Expressed as Equivalence of Two Functions
- Second function described with a LAMBDA expression

Example Using Trig Library (cont.)

```
sp4 :

|-----
{1}   FORALL (p, v: posreal): syst(v, p) = (LAMBDA x: 1 - sq(sin(x)))

Rule? (skosimp*)

|-----
{1}   syst(v!1, p!1) = (LAMBDA x: 1 - sq(sin(x)))

Rule? (apply-extensionality 1 :hide? t)           or TAB E

|-----
{1}   syst(v!1, p!1)(x!1) = 1 - sq(sin(x!1))

Rule? (lemma "sp3")

{-1}  FORALL (p, v: posreal, x: nnreal): syst(v, p)(x) = 1 - sq(sin(x))
|-----
[1]   syst(v!1, p!1)(x!1) = 1 - sq(sin(x!1))

Rule? (inst?)
x: nnreal gets x!1, p: posreal gets p!1, v: posreal gets v!1,
```

Q.E.D.

Analysis Library

Limit of a functions [T -> real] at a point

lim_of_functions, lim_of_composition

Limits and operations on sequences of reals

convergence_ops, convergence_sequences

Continuous functions [T -> real]

continuous_functions, composition_continuous,
continuous_functions_props,
unif_cont_fun,
inverse_continuous_functions, continuous_linear

Differential Calculus

derivatives, derivative_props, chain_rule,
derivatives_more, sqrt_derivative, derivative_inverse

Integral Calculus

integral_def, integral_cont, integral_split ,
integral, fundamental_theorem,
table_of_integrals, integral_chg_var, integral_diff_doms

What Does Continuity Look Like?

```
continuous(f, x0) IFF  
FORALL epsilon : EXISTS delta : FORALL x :  
abs(x - x0) < delta IMPLIES abs(f(x) - f(x0)) < epsilon
```

```
sum_continuous : LEMMA continuous(f1, x0) AND continuous(f2, x0)  
IMPLIES continuous(f1 + f2, x0)
```

```
diff_continuous : LEMMA continuous(f1, x0) AND continuous(f2, x0)  
IMPLIES continuous(f1 - f2, x0)
```

```
prod_continuous : LEMMA continuous(f1, x0) AND continuous(f2, x0)  
IMPLIES continuous(f1 * f2, x0)
```

```
const_continuous: LEMMA continuous(const_fun(u), x0)
```

```
scal_continuous : LEMMA continuous(f, x0)  
IMPLIES continuous(u * f, x0)
```

Derivatives

How do you define a subtype of the reals suitable for defining derivatives?

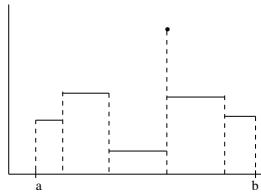
```
derivatives[T: TYPE FROM real]: THEORY
ASSUMING

connected_domain : ASSUMPTION
  FORALL (x, y : T), (z : real) :
    x <= z AND z <= y IMPLIES T_pred(z)

not_one_element : ASSUMPTION
  FORALL (x : T) : EXISTS (y : T) : x /= y

ENDASSUMING
```

Integrals



```
integral?(a:T,b:{x:T|a<x},f:[T->real],S:real): bool =
  (FORALL (epsi: posreal): (EXISTS (delta: posreal):
    (FORALL (P: partition(a,b)):
      width(a,b,P) < delta IMPLIES
        (FORALL (R: (Riemann_sum?(a,b,P,f)))):
          abs(S - R) < epsi))))
```

From this definition we can construct a predicate integrable? and a function integral which is defined on integrable? functions:

```
integrable?(a:T,b:{x:T|a<x},f:[T->real]): bool =
  (EXISTS (S: real): integral?(a,b,f,S))

integral(a:T,b:{x:T|a<x}, ff: { f | integrable?(a,b,f)}):
  {S: real | integral?(a,b,ff,S)}
```

Integral Theorems

It is expected that most users of this formalization of the integral, will only need the theorems in two PVS theories: integral and fundamental_theorem. Quick reference guide:

Theorem	Lemma Name
$\int_a^a f(x) dx = 0$	Integral_a_to_a
$\int_a^b c dx = c * (b - a)$	Integral_const_fun
$\int_b^a f(x) dx = - \int_a^b f(x) dx$	Integral_rev
$\int_a^b cf(x) dx = c \int_a^b f(x) dx$	Integral_scal
$\int_a^b f(x) + g(x) dx = \int_a^b f(x) dx + \int_a^b g(x) dx$	Integral_sum
$\int_a^b f(x) - g(x) dx = \int_a^b f(x) dx - \int_a^b g(x) dx$	Integral_diff
$\int_a^b f \text{ WITH } [(y) := yv] dx = \int_a^b f dx$	Integral_chg_one_pt
$f(x) \geq 0 \supset \int_a^b f dx \geq 0$	Integral_ge_0
$ f(x) < M \supset \int_a^b f(x) dx \leq M * (b - a)$	Integral_bounded
$f \text{ integrable} \supset f(x) < B$	Integrable_bounded
$f \text{ continuous} \supset f \text{ integrable}$	continuous_Integrable?
$\int_a^b f(x) dx + \int_b^c f(x) dx \int_a^c f(x) dx$	Integral_split
For step function $f : \int_a^b f(x) dx = \sum_{i=1}^n c_i(x_i - x_{i-1})$	step_function_on_integral
$F(x) = \int_a^x f dt \supset F' = f$	fundamental1
$\int_a^b f(t) dt = F(b) - F(a)$	fundamental3

The Inexp Library

There are two versions: lnexp and lnexp_fnd. The first is axiomatic and the second is foundational. But from a user's viewpoint they should be interchangeable.

```

ln_exp           -- definitions
expt            -- a^x where x real and a >= 0
hyperbolic       -- e.g. sinh cosh functions
exp_series      -- exp as infinite series
ln_series       -- ln as taylor series
exp_approx      -- exp approximation
ln_approx       -- ln approximation
ln_exp_series_alt -- alternate ln/exp series formulation
ln_exp_ineq     -- inequalities involving ln/exp

```

The Inexp_fnd Library

In the foundational version:

```
ln(x: posreal): real = Integral(1,x, (LAMBDA (t: posreal): 1/t))

ln_derivable    : LEMMA derivable(ln) AND
                  deriv(ln) = (LAMBDA (t: posreal): 1/t)

ln_continuous   : LEMMA continuous(ln)

ln_1            : LEMMA ln(1) = 0

ln_mult         : LEMMA ln(px*py) = ln(px) + ln(py)
ln_div          : LEMMA ln(px/py) = ln(px) - ln(py)
ln_expt        : LEMMA ln(px^i) = i*ln(px)
```

The Inexp Library (cont.)

In the axiomatic version:

```
ln(x: posreal): real % = Integral(1,x, (LAMBDA (t: posreal): 1/t)

%  ln_derivable    : LEMMA derivable(ln) AND
%                      deriv(ln) = (LAMBDA (t: posreal): 1/t)

%  ln_continuous   : LEMMA continuous(ln)

ln_1            : AXIOM ln(1) = 0

ln_mult         : AXIOM ln(px*py) = ln(px) + ln(py)
ln_div          : LEMMA ln(px/py) = ln(px) - ln(py)
ln_expt        : LEMMA ln(px^i) = i*ln(px)
```

**Notice that the “analysis” stuff is commented out.
Should we import an axiomatic version of “analysis” instead?**

Series Library

```
series           -- definition + properties of infinite series
power_series    -- definition + properties of power series
trig_fun         -- definition of trigonometric functions
trig_props       -- basic trig identities
nth_derivatives -- nth derivative of a function
taylors          -- Taylor's theorem
taylor_series    -- expansion into Taylor's series
power_series_deriv -- derivative of a power series
power_series_integ -- integral of a power series
```

Series Library

```
series(a,m) =  (LAMBDA n:   \      a(j) )
                  /
                  ----
                  j = m

series(a) : sequence[real] = (LAMBDA (n: nat): sigma(0, n, a))
series(a,m): sequence[real] = LAMBDA (n: nat): sigma(m, n, a))

conv_series?(a,m): bool = convergent(series(a,m))

inf_sum(m, a: {s | conv_series?(s,m)}): real = limit(series(a,m))

powerseq(a,x): sequence[real] = (LAMBDA (k: nat): a(k)*x^k)

powerseries(a)(x): sequence[real] = series(powerseq(a,x))
```

Auto Rewrites in Libraries

- Throughout the NASA libraries you will find statements like

```
AUTO_REWRITE+ dbl_to_pair_inverse_1    % dbl(dbl_to_pair(D)) =  
AUTO_REWRITE+ Riemann?_Rie  
AUTO_REWRITE+ sqrt_0
```

- Why have we added these?

- ANSWER: They can remove some of the tedium:

$$\sqrt(0) * \cos(2\pi + x!1) = \sqrt(0) / \ln(x!1 * y!1)$$

- Certain simplifications are done automatically:

$$\sqrt(0) \rightarrow 0$$
$$\sqrt(0) \rightarrow 0$$

Automatic Rewriting

- Idea is simple, but powerful

- Start with a lemma that is an equality, e.g.

$$\text{sqrt_square: LEMMA } \sqrt(nnx * nnx) = nnx$$

- If the left hand side matches a sub-expression in a formula, e.g.

$$\sqrt(\exp(x) * \exp(x)) * \ln(1 - y * y)$$

- then replace with right hand side:

$$\exp(x) * \ln(1 - y * y)$$

Several Ways To Turn On Auto-Rewrites

- Put AUTO_REWRITE+ sqrt_square in your theory somewhere
- Issue a command like (auto-rewrite "sqrt_square")
- Issue a command like (auto-rewrite-theory "sqrt")⁴
- Use a strategy that turns on auto-rewrites, e.g. (realprops) or (grind)

⁴This should only be done with theories specifically designed for auto-rewriting.

Automatic Rewriting (cont.)

For recursively defined functions

```
factorial(n): recursive posnat =
  (if n = 0 then 1 else n*factorial(n-1) endif) MEASURE n
```

the result can be dramatic:

```
|-----
[1] factorial(12) = ...
```

```
Rule? (AUTO-REWRITE "factorial")
Rule? (ASSERT)
```

```
factorial rewrites factorial(0) to 1
factorial rewrites factorial(1) to 1
factorial rewrites factorial(2) to 2
factorial rewrites factorial(3) to 6
factorial rewrites factorial(4) to 24
factorial rewrites factorial(5) to 120
factorial rewrites factorial(6) to 720
factorial rewrites factorial(7) to 5040
factorial rewrites factorial(8) to 40320
factorial rewrites factorial(9) to 362880
factorial rewrites factorial(10) to 3628800
factorial rewrites factorial(11) to 39916800
factorial rewrites factorial(12) to 479001600
```

What About?

What happens if you do the following?

```
u,v: VAR Vector
```

```
add_comm: LEMMA u+v = v+u
```

```
AUTO_REWRITE+ add_comm
```

Some AUTO-REWRITES In the reals Library

```
AUTO_REWRITE- sq_neg_minus
AUTO_REWRITE+ sq_0
AUTO_REWRITE+ sq_1
AUTO_REWRITE+ sq_abs
AUTO_REWRITE+ sq_abs_neg
AUTO_REWRITE+ sq_neg          % LEMMA sq(-a) = sq(a)
AUTO_REWRITE+ sq_0            % LEMMA sq(0) = 0
AUTO_REWRITE+ sq_1            % LEMMA sq(1) = 1
AUTO_REWRITE+ sq_abs          % LEMMA sq(abs(a)) = sq(a)
AUTO_REWRITE+ sq_abs_neg      % LEMMA sq(abs(-a)) = sq(a)
AUTO_REWRITE+ sqrt_4
AUTO_REWRITE+ sqrt_9
AUTO_REWRITE+ sqrt_16
AUTO_REWRITE+ sqrt_25
AUTO_REWRITE+ sqrt_0
AUTO_REWRITE+ sqrt_1
AUTO_REWRITE+ sqrt_square
AUTO_REWRITE+ sqrt_sq
AUTO_REWRITE+ sqrt_sq_neg
AUTO_REWRITE+ sq_sqrt
```

Using “auto-rewrite-theory”

- You can turn on all of the rewrites in a theory using command “auto-rewrite-theory”, e.g.

```
(auto-rewrite-theory "real_props")
```

- Another way is to create a theory with a bunch of AUTO_REWRITE+ commands in it
- If you import this theory, all of the rewrites are turned on. For example:

```
sqrt_rew: THEORY
%-----
%
% Square root AUTO-REWRITE+ definitions to increase automation
%
%-----
```

BEGIN

```
IMPORTING sqrt, sq_rew
```

```
nnx, nny, nnz, nna : VAR nonneg_real
x,y,z,xx: VAR real
```

```
AUTO_REWRITE+ sqrt_0          % : LEMMA sqrt(0) = 0
AUTO_REWRITE+ sqrt_1          % : LEMMA sqrt(1) = 1
```

```

AUTO_REWRITE+ sqrt_eq_0      % : LEMMA sqrt(nnx) = 0 IMPLIES nnx = 0
AUTO_REWRITE+ sqrt_def       % : LEMMA sqrt(nnx) * sqrt(nnx) = nnx
AUTO_REWRITE+ sqrt_square    % : LEMMA sqrt(nnx * nnx) = nnx

AUTO_REWRITE+ sqrt_sq        % : LEMMA x >= 0 IMPLIES sqrt(sq(x)) = x
AUTO_REWRITE+ sqrt_sq_neg    % : LEMMA x < 0 IMPLIES sqrt(sq(x)) = -x
AUTO_REWRITE+ sqrt_sq_abs    % : LEMMA sqrt(sq(x)) = abs(x)

sqrt_4                  : LEMMA sqrt(4) = 2
sqrt_9                  : LEMMA sqrt(9) = 3
sqrt_16                 : LEMMA sqrt(16) = 4
sqrt_25                 : LEMMA sqrt(25) = 5
sqrt_36                 : LEMMA sqrt(36) = 6
sqrt_49                 : LEMMA sqrt(49) = 7
sqrt_64                 : LEMMA sqrt(64) = 8
sqrt_81                 : LEMMA sqrt(81) = 9
sqrt_100                : LEMMA sqrt(100) = 10

AUTO_REWRITE+ sqrt_4
AUTO_REWRITE+ sqrt_9
AUTO_REWRITE+ sqrt_16
AUTO_REWRITE+ sqrt_25
AUTO_REWRITE+ sqrt_36
AUTO_REWRITE+ sqrt_49
AUTO_REWRITE+ sqrt_64
AUTO_REWRITE+ sqrt_81
AUTO_REWRITE+ sqrt_100

sqrt_factor_left  : LEMMA sqrt(nna * nnx * nnx) = sqrt(nna) * nnx
sqrt_factor_middle: LEMMA sqrt(nnx * nna * nnx) = sqrt(nna) * nnx
sqrt_factor_right : LEMMA sqrt(nnx * nnx * nna) = sqrt(nna) * nnx

AUTO_REWRITE+ sqrt_factor_left
AUTO_REWRITE+ sqrt_factor_middle
AUTO_REWRITE+ sqrt_factor_right

% sqrt_times_rev  : LEMMA sqrt(nny) * sqrt(nnz) = sqrt(nny * nnz)
% sqrt_div_rev    : LEMMA nnz /= 0 IMPLIES
%                      sqrt(nny) / sqrt(nnz) = sqrt(nny / nnz)

% AUTO_REWRITE+ sqrt_times_rev
% AUTO_REWRITE+ sqrt_div_rev

END sqrt_rew

```

Structures Library

```
% ----- arrays -----  
top_array  
min_array      % defines min function over an array  
max_array      % defines max function over an array  
permutations   % permutations defined using arrays  
sort_array     % defines a sort function over arrays  
sort_array_lems % relationship between sort and min and max  
array_ops       % array operations  
majority_array  % defines majority function over an array  
  
% ----- sequences -----  
top_seq  
max_seq        % defines max function over an sequence  
min_seq        % defines min function over a sequence  
permutations_seq % permutations defined using arrays  
majority_seq    % defines majority function over finite seq  
bubblesort      % bubble sort correctness theorem  
sort_seq        % defines a sort function over sequences  
sort_seq_lems   % relationship between sort and min and max  
seq2set         % convert sequence to a set
```

Structures Library (cont.)

```
% ----- bags -----  
top_bags  
bags          % fundamental definitions and properties  
bags_aux      % definition of filter rest and choose  
bags_to_sets  % converts bags to sets  
finite_bags   % basic definitions and lemmas  
finite_bags_lems % lemmas need induction  
finite_bags_aux % lemmas about filter rest and choose  
finite_bags inductions % induction schemes  
bag_filters    % filtering linearly-ordered bags pigeonhole  
                % majority pigeonhole results and overlap  
                % existence proof  
majority_vote  % defines a majority vote over finite bags  
middle_value_select % defines middle value selection over finite  
fault_masking_vote % proves an equivalence between majority and  
                    % value selection  
finite_bags_minmax % defines the minimum and maximum of a finite  
                    % ordered bag.
```

Vectors Library

```
vectors           % N-dimensional vectors and operations
vect2D            % Define 2-D Vector from N-dimensional vectors
vect3D            % Define 3-D Vector from N-dimensional vectors

vectors_cos       % Law of cosines for n-D vectors
vectors2D_cos    % Law of cosines for 2D vectors
vectors3D_cos    % Law of cosines for 3D vectors

position .. 2D ..3D    % using vectors for position, dist func

lines, lines2D, lines3D % Using vectors to define lines, motion

law_cos_pos2D, .. 3D    % Law of cosines

closest_approach, _2D/3D % closest point in time
perpendicular2D, ..3D    % line perpendicular to line through pt
vectors_sign2D          % signs of vector dot product
intersections2D         % finding intersection points of lines
matrices               % Theory of matrices
```

Some Auto-Rewrites in the Vectors Library

```
AUTO_REWRITE+ add_zero_left           % zero + v = -v
AUTO_REWRITE+ add_zero_right          % v + zero = v
AUTO_REWRITE+ sub_zero_left           % zero - v = -v
AUTO_REWRITE+ sub_zero_right          % v - zero = v
AUTO_REWRITE+ sub_eq_zero             % u - v = zero IFF u = v
AUTO_REWRITE+ sub_eq_args             % v - v = zero
AUTO_REWRITE+ neg_add_left            % -v + v = zero
AUTO_REWRITE+ neg_add_right           % v + -v = zero
AUTO_REWRITE+ dot_zero_left           % zero * v = 0
AUTO_REWRITE+ dot_zero_right          % v * zero = 0
AUTO_REWRITE+ scal_zero              % a * zero = zero
AUTO_REWRITE+ scal_0                 % 0 * v = zero
AUTO_REWRITE+ sqv_zero                % sqv(zero) = 0
AUTO_REWRITE+ add_neg_sub             % v + -u = v - u
AUTO_REWRITE+ neg_neg                % --v = v
AUTO_REWRITE+ dot_scal_left           % (a*u)*v = a*(u*v)
AUTO_REWRITE+ dot_scal_right          % u*(a*v) = a*(u*v)
AUTO_REWRITE+ dot_scal_assoc          % a*(b*u) = (a*b)*u
AUTO_REWRITE+ dot_scal_canon          % (a*u)*(b*v) = (a*b)*(u*v)
AUTO_REWRITE+ sqv_neg                 % sqv(-v) = sqv(v)
AUTO_REWRITE+ sqrt_sqv_sq              % sqrt(sqv(v)) * sqrt(sqv(v)) = sqv(v)
AUTO_REWRITE+ norm_neg                % norm(-u) = norm(u)
AUTO_REWRITE+ comp_zero_x              % zero'x = 0
AUTO_REWRITE+ comp_zero_y              % zero'y = 0
AUTO_REWRITE+ add_cancel              % v + w - v = w
AUTO_REWRITE+ sub_cancel              % v - w - v = -w
AUTO_REWRITE+ add_cancel_neg           % -v + w + v = w
AUTO_REWRITE+ sub_cancel_neg           % -v - w + v = -w
AUTO_REWRITE+ add_cancel2              % w - v + v = w
AUTO_REWRITE+ add_cancel_neg2           % w + v - v = w
```

```

AUTO_REWRITE-    zero
AUTO_REWRITE-    add_comm      % u+v = v+u
AUTO_REWRITE-    dot_comm      % u*v = v*u
AUTO_REWRITE-    dot_assoc     % a*(v*w) = (a*v)*w
AUTO_REWRITE-    sqv_lem       % sqv(u) = u*u
AUTO_REWRITE-    sqv_sym       % sqv(u-v) = sqv(v-u)
AUTO_REWRITE-    norm_sym      % norm(u-v) = norm(v-u)
AUTO_REWRITE-    dot_sq_norm   % u*u = sq(norm(u))

```

Floating Point (float) Library

High-level model (SB)

float, axpy,

Hardware-level (PM)

```

IEEE_854,
IEEE_854_defs,
infinity_arithmetic,
comparison1,
NaN_ops,
arithmetic_ops,
IEEE_854_remainder,
IEEE_854_fp_int,
real_to_fp, over_under,
IEEE_854_values,
round, fp_round_aux,
sum_lemmas,
enumerated_type_defs, sum_hack,

```

Equivalence between the two models (SB)

IEEE_link

Complex Numbers (complex) Library

```
complex_types      % Basic Definitions of Complex Number Types
polar              % Polar coordinate complex numbers
arithmetic         % Basic Arithmetic on Complex Numbers
exp                % Complex logarithm and exponential functions

complex: TYPE+ FROM number_field

i : complex
i_axiom:           AXIOM i*i = -1

nf:                 VAR numfield
x,x0,x1,y,y0,y1: VAR real

complex_characterization: AXIOM complex_pred(nf)
                           IFF EXISTS x,y: nf = x+y*i

real_complex:       LEMMA complex_pred(x)
i_not_real:         LEMMA r /= i
```

Algebra Library

Groups:

groupoid,	finite_groupoid,
commutative_groupoid,	finite_commutative_groupoid,
monad,	finite_monad,
commutative_monad,	finite_commutative_monad,
semigroup,	finite_semigroup,
commutative_semigroup,	finite_commutative_semigroup,
monoid,	finite_monoid,
commutative_monoid,	finite_commutative_monoid,
cyclic_monoid,	
group,	finite_group,
abelian_group,	finite_abelian_group,
symmetric_groups,	
group_test	

Algebra Library (cont.)

Ring/Field-like Mathematical Structures

```
%      [T:Type+,+,*: [T,T->T] ,zero:T]   or   [T:Type+,+,*: [T,T->T] ,zero:T]
%
%-----+-----+
%          File           | Mathematical Structure of *
%-----+-----+
```

ring,	% semigroup[T,*]
commutative_ring,	% commutative_semigroup[T,*]
ring_nz_closed,	% semigroup[T,*], semigroup[nz_T,*]
ring_with_one,	% monoid[T,*,one]
commutative_ring_with_one,	% commutative_monoid[T,*,one]
integral_domain,	% commutative_semigroup[T,*]
	% commutative_semigroup[nz_T,*]
division_ring,	% monoid[T,*,one], group[nz_T,*,one]
field	% commutative_monoid[T,*,one]
	% abelian_group[nz_T,*,one]

Graph Theory Library (graphs)

Main theories:

circuits	-- theory of circuits
graphs	-- fundamental definition of a graph
graph_connected	-- all connected defs are equivalent
graph_deg	-- definition of degree
graph_inductions	-- vertex and edge inductions for graphs
graph_ops	-- delete vertex and delete edge operations
max_subgraphs	-- maximal subgraphs with specified property
max_subtrees	-- maximal subtrees with specified property
menger	-- menger's theorem
min_walks	-- minimum walk satisfying a property
paths	-- fundamental def and properties about paths
ramsey_new	-- Ramsey's theorem
subgraphs	-- generation of subgraphs from vertex sets
subgraphs_from_walk	-- generation of subgraphs from walks
subtrees	-- subtrees of a graph
tree_circ	-- theorem that tree has no circuits
tree_paths	-- theorem tree has only one path between vertices
trees	-- fundamental definition of trees
walk_inductions	-- induction on length of a walk
walks	-- fundamental definition and properties of walks

Graph Theory Library (graphs)

```
dbl(x,y): set[T] = {t: T | t = x OR t = y}

doubleton: TYPE = {S: set[T] | EXISTS x,y: x /= y
                    AND S = dbl(x,y) }

pregraph: TYPE = [# vert : finite_set[T] ,
                  edges: finite_set[doubleton[T]] #]

graph: TYPE = {g: pregraph |
               (FORALL (e: doubleton[T]): edges(g)(e)
                IMPLIES (FORALL (x:T): e(x) IMPLIES vert(g)(x)))}
```

DiGraph Theory Library (digraphs)

```
edgetype: TYPE = pair[T]

edg(x, (y:{y:T | y /= x})): edgetype = (x,y)

predigraph: TYPE = [# vert : finite_set[T] ,
                     edges: finite_set[edgetype] #]

e: VAR edgetype

directed_graph: TYPE = {g: predigraph | FORALL e: edges(g)(e)
                           IMPLIES LET (x,y) = e IN
                           vert(g)(x) AND vert(g)(y) }

digraph : TYPE = directed_graph
```

The sets_aux Library

```
power_sets          % cardinality and finiteness

card_comp           % compare the cardinality of any two types
card_finite         % card_comp vs. card(S) from finite_sets
card_power          % card[T] < card[set[T]]
card_power_set      % card(S) < card(powerset(S))
card_sets_lemmas    % relationships of set operations to cardinal
card_single         % single type properties of card_comp_set pre
countability        % definition of (un)countable sets
countable_props     % properties of (un)countable sets
countable_set       % some countable sets of numbers (e.g. integer)
countable_setofsets % operations on countable families of sets
countable_types     % countability of some prelude types
infinite_card       % card_comp implications for finiteness
infinite_image      % infinite images of a set under some function
infinite_sets       % cardinality of infinite set add and remove
```

Orders Library

```
bounded_orders      % definitions of lub, glb, (complete) lattices
closure_ops          % reflexive, symmetric, transitive, etc. closure
complementary_lattices % lattices with a "complement" function
complete_lattices    % every set is tightly bounded
finite_orders         % properties of an order on a finite type
fixed_points          % fixed points characterized by prefixed point
lattices              % operations that preserve tight-boundedness
lower_semilattices    % definition of binary glb function
new_mu calculus_prop % a simulation of fixedpoints@mu calculus_property
pointwise_orders      % lifting an order to functions
total_lattices        % a lattice defined by a total order
chain                 % totally ordered subsets of a poset
chain_chain            % chains of chains in inclusion order
converse_zorn          % lower bound on all chains => min. element
isomorphism             % isomorphisms between ordered sets
kuratowski             % there exists a maximal chain in any set
order_strength          % strengthenings and weakenings of orders
set_dichotomous        % injective map exists between any 2 sets
well_ordering           % every set has a well-ordering relation
zorn                   % upper bound on all chains => max. element
infinite_pigeonhole    % [infinite_domain -> finite_range] enumeration
```

Probability Library

```
IMPORTING continuous_functions_aux,  
        probability_measure,  
        probability_space,  
        conditional,  
        expectation
```

ACCoRD Library

- Horizontal Criteria
- Vertical Criteria
- Loss of Separation
- 2-D Conflict Detection (cd2d)
- 2-D Bands
- 3-D Criteria
- 3-D Horizontal Resolution Algorithms
- Vertical Resolution Algoritms
- 3-D Conflict Detection (cd3d)
- Conflict Resolution (cr3d)
- 3-D Bands
- Other Algorithms
- Flight plans
- Conflict detection for state-based ownship and intent-based intr
- Conflict detection for intent-based ownship and intruder (cd3d_i
- Prevention bands for state-based ownship and intent-based intrud

Fault Tolerance Library

```
majority_integration,  
exact_reduce_integration,  
reduce_synch,  
inexact_reduce,  
convergence_top
```

```
virtual_clock_top
```

Naming Conventions

- Lemmas should begin with the function name.
- Key defining property should be labeled `_def`.
- Common useful rewrite label it `_rew`.
- Common alternate or simpler version label it `_lem`.

abbrev	meaning
<code>_0</code>	value of function at 0
<code>_eq_0</code>	function equals 0: $f(x) = 0$ IFF ...
<code>_eq_args</code>	$f(a,a) = \dots$
<code>_neg</code>	value of function for negated argument $f(-x)$
<code>_plus</code>	value of function for sum of arguments $f(x+y)$
<code>_plus1</code>	value of function for $f(x+1)$
<code>_minus</code>	value of function for difference of arguments $f(x-y)$
<code>_disj</code>	disjoint
<code>_dist</code>	distributive
<code>_comm</code>	commutative: $f(a,b) = f(b,a)$
<code>_assoc</code>	associative: $f(a,f(b,c)) = f(f(a,b),c)$
<code>_sym</code>	symmetry: $f(-a) = f(a)$
<code>_incr</code>	$f(a) \leq f(b)$ IFF $a \leq b$
<code>_decr</code>	$f(a) \geq f(b)$ IFF $a \leq b$
<code>_strict_incr</code>	$f(a) < f(b)$ IFF $a < b$
<code>_strict_decr</code>	$f(a) > f(b)$ IFF $a < b$
<code>_fix_pt</code>	value of the defined function is a fixed point
<code>_card</code>	cardinality value
<code>_lb</code>	lower bound
<code>_ub</code>	upper bound